Public Key Linting for ML-KEM and ML-DSA

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Motivation

Incident Dashboard





partSIGN: Findings in 2025 FTSI

Audit - Audit Incident Report #5 -

1965808 ASSIGNED

- certification authorities (CA) issue certificates used to verify the identity of entities
 - crucial for security and functioning of protocols like TLS, S/MIME, ...
 - essential to establish trust among users
- several requirements regarding certification exist
 - \rightarrow failing to comply with them leads to incidents

Incident Dashboard

GlobalSign: Non-BR-Compliant Certificate Issuance RSA key smaller than 2048 bits	1393557	RESOLVED
Sectigo: Certificates with RSA keys where modulus is not divisible by 8	1653504	RESOLVED
GDCA: Misissuance of certificates with small RSA keys	1467414	RESOLVED

- incidents harm a CA's reputation and more generally public key infrastructures
- many incidents are related to the content of a certificate
- can be avoided, if proper mechanisms are in place

Linting

- linting: process of analyzing the content of a certificate w.r.t. predefined rules ("lints")
- examples for RSA lints:
 - the length of the modulus is one of the specified values in FIPS 186-3
 - the modulus and the public exponent are odd numbers
 - the modulus is not a power of a prime
 - the modulus has no factors smaller than 752

Lint Example

Lint: The modulus has no factors smaller than 752

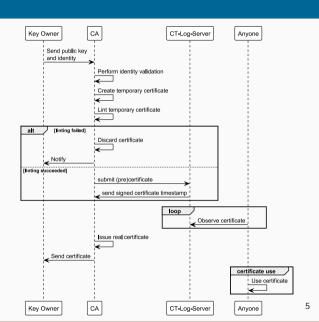
Check that gcd(n, r) = 1 where

 $r = 1451887755777639901511587432083070202422614380984889313550570919 \\ 6593151770659565743590789126541491676439926842369913057775743308 \\ 3166651158914570105971074227669275788291575622090199821297575654 \\ 3223550490431013061082131040808010565293748926901442915057819663 \\ 730454818359472391642885328171302299245556663073719855 \, ,$

is the product of the 132 primes from 3 to 751.

Linting Process

- CA creates temporary certificate and starts the linting process
- certificate is checked for each lint in the linting library
- if one lint fails, the entire linting process fails → the CA discards the certificate and notifies the key owner
- otherwise: certificate is sent to a certificate transparency log server, finalized, and issued



Current State of Linting Regulations

Ballot SC-75 of CA/Browser Forum...

- ...renders pre-sign linting mandatory from March 2025
- ...states that validating the key material is a responsibility of the CA

 $\underline{\text{Currently:}} \text{ these requirements cover RSA and elliptic curve keys}$

Linting Post-Quantum Schemes

- advancing standardization of post-quantum cryptographic (PQC) schemes:
 - finalized standards: ML-KEM, ML-DSA, SLH-DSA
 - also selected for standardization: Falcon, HQC
 - ongoing NIST process for additional digital signature schemes
- need for preparing the IT security infrastructure for integrating post-quantum schemes
- linting has not been studied for PQC schemes!

Goal: initiate the study of linting for PQC schemes by analyzing the public keys of ML-KEM and ML-DSA

Background: ML-KEM

ML-KEM: General Information

- NIST standard for key-encapsulation mechanisms (KEM) using module lattices
- based on the KEM CRYSTALS-Kyber
- high-level construction:

public-key encryption scheme: K-PKE FO-transform KEM: ML-KEM

- distinction between external and internal components:
 - external: generate randomness, check whether randomness generation was successful, and call their internal counterparts
 - internal: actual steps of the procedures ← here the public key is used!

ML-KEM: Key Generation

Algorithm ML-KEM.KeyGen_internal(d, z)

Input randomness $d, z \in \mathbb{B}^{32}$

Output $ek \in \mathbb{B}^{384k+32}$, $dk \in \mathbb{B}^{768k+96}$

- 1: $(ek_{PKE}, dk_{PKE}) \leftarrow K-PKE.KeyGen(d)$
- 2: $ek \leftarrow ek_{PKE}$
- 3: $dk \leftarrow (dk_{PKE}||ek||H(ek)||z)$
- 4: **return** (*ek*, *dk*)

ML-KEM: Key Generation of the underlying PKE

Algorithm K-PKE.KeyGen(d)

```
Input: randomness d \in \mathbb{B}^{32}
```

Output:
$$ek_{\mathsf{PKE}} \in \mathbb{B}^{384k+32}$$
, $dk_{\mathsf{PKE}} \in \mathbb{B}^{384k}$

1:
$$(\rho, \sigma) \leftarrow G(d||k)$$

2:
$$N \leftarrow 0$$

3: **for**
$$(i \leftarrow 0; i < k; i++)$$
 do

4: **for**
$$(j \leftarrow 0; j < k; j++)$$
 do

5:
$$\hat{A}[i,j] \leftarrow \mathsf{SampleNTT}(\rho||j||i)$$

6: end for

7: end for

8: **for**
$$(i \leftarrow 0; i < k; i++)$$
 do

9:
$$s[i] \leftarrow \mathsf{SamplePolyCBD}_{\eta_1}(\mathsf{PRF}_{\eta_1}(\sigma, N))$$

10:
$$N \leftarrow N + 1$$

11: end for

Algorithm K-PKE.KeyGen(d)

12: **for**
$$(i \leftarrow 0; i < k; i++)$$
 do

13:
$$e[i] \leftarrow \mathsf{SamplePolyCBD}_{\eta_1}(\mathsf{PRF}_{\eta_1}(\sigma, N))$$

14:
$$N \leftarrow N + 1$$

16:
$$\hat{s} \leftarrow \mathsf{NTT}(s)$$

17:
$$\hat{e} \leftarrow \mathsf{NTT}(e)$$

18:
$$\hat{t} \leftarrow \hat{A} \circ \hat{s} + \hat{e}$$

19:
$$ek_{PKE} \leftarrow ByteEncode_{12}(\hat{t}) \| \rho$$

20:
$$dk_{PKE} \leftarrow ByteEncode_{12}(\hat{s})$$

21: **return** (
$$ek_{PKE}$$
, dk_{PKE})

Methodology

Finding Lints

- consider properties of the certificate and the public key itself
- check whether all rules from the standards are fulfilled and if the properties of an honestly generated key are given
- take into account the input validation checks described in the standards

Lint Classes

We introduce 5 lint classes:

- **INTER:** interoperability lints focus on certificate properties which assist applications in properly communicating with each other
- DIM: dimension lints test whether the size of certain objects is correct
- DOM: domain lints test whether the "type" of the objects is correct
- DIS: distribution lints verify distribution properties of objects
- GEN: lints that work generically for any scheme

Completeness of Lints

- lint classes: new ones might need to be introduced for future lints!
- dimension and domain lints: complete
- <u>distribution lints:</u> incomplete
- interoperability and generic lints: incomplete
 - ightarrow only starting point for post-quantum linting
 - \rightarrow more lints can be added over time using our proposed formal description

Application Lints

- applications (e.g., email client, browser) extract the pk from the certificate in order to use it
- in the case of ML-KEM and ML-DSA, the application must perform further operations, like expanding the key
 - ightarrow the expanded pk depends on the implementation of the expanding algorithm used by the application
 - \rightarrow lints for the expanded pk must be performed at the application side (if at all)

Implementation

- our lints are implemented in Java and Rust
 - lints which examine properties of the certificate \rightarrow Java (using methods provided by BouncyCastle)
 - lints that check the correctness of the key within the certificate \rightarrow Rust
- Rust allows us to enforce the size of inputs on the type level, such that dimension lints get covered by the API
- we give test vectors for keys with and without errors

ML-KEM Lints

Overview

Lint	Classification	Identifier
key usage	interoperability	INTER_01
pk aid enc	interoperability	INTER_02
ek length	dimension	DIM_01
ek seed length	dimension	DIM_02
ek matrix dimension	dimension	DIM_03
ek vector dimension	dimension	DIM_04
ek matrix entries	domain	DOM_01
ek vector entries	domain	DOM_02
ek seed entry frequency	distribution	DIS_01
ek seed entry run	distribution	DIS_02
ek seed small/large entries	distribution	DIS_03
ek matrix entry frequency	distribution	DIS_04
ek matrix entry run	distribution	DIS_05
ek matrix small/large entries	distribution	DIS_06
known enc key	generic	GEN_01
ek encoding	generic	GEN_02

Example: Interoperability Lint

ML_KEM_INTER_01

- concerns the key usage extension in a certificate, which specifies for which types of use the certificate's public key can be used
- lint checks that the value of the key usage extension is compatible with the ML-KEM algorithm
- only values related to key encryption are compatible

Example: Dimension Lint

ML_KEM_DIM_01

- based on one of the tests for input validation described in the ML-KEM standard
- lint checks the length of the encoded encapsulation key ek
- the correct length is 384k+32 bytes for $k\in\{2,3,4\}$ depending on the ML-KEM parameter set

Example: Distribution Lint

ML_KEM_DIS_01

- checks if the seed ρ contains the same element an amount of times that is unlikely for a pseudorandomly sampled value from \mathbb{B}^{32}
- outputs an error if there are at least x = 20 occurrences of the same byte

Conclusion

Conclusion

- we initiate the study of PQC linting and provide a framework
- challenges:
 - for PQC schemes, security is often related to certain elements "looking random" or following a pre-defined distribution
 - usage of seeds in ML-KEM and ML-DSA: properties of the expanded public key cannot be checked by the CA (depends on the implementation of the expansion algorithm)
 - \rightarrow testing of the implementations is necessary
- only a starting point: linting is very scheme-specific!